

Optimal Dispatch of Multiple Photovoltaic Integrated 5G Base Stations for Active Distribution Network Demand Response

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Multiple 5G base stations (BSs) equipped with distributed photovoltaic (PV) generation devices and energy storage (ES) units participate in active distribution network (ADN)

demand response (DR), which is expected to be the best way to reduce the energy cost of 5G BSs and provide flexibility resources for the ADN. However, the standalone PVintegrated 5G BS has the characteristics of wide distribution, small volume, and large load fluctuations, which will bring strong uncertainty to the ADN by directly participating in the DR. Therefore, a system architecture for multiple PV-integrated 5G BSs to participate in the DR is proposed, where an energy aggregator is introduced to effectively aggregate the PV energy and ES resources of 5G BSs. Then, a two-stage optimal dispatch method is proposed. Specifically, in the large-timescale DR planning stage, an incentive mechanism for multiple PV-integrated 5G BSs to participate in the DR is constructed based on the contract theory, which ensures that multiple 5G BSs respond to and satisfy the peakshaving demand of the ADN. In the small-timescale online energy optimization stage, based on the energy sharing mode among 5G BSs, a Lyapunov-based online energy optimization algorithm is proposed to optimize the shared energy flow between the internal layer and the interactive layer of 5G BSs, which further improves PV absorption and ensures operation stability of ES in the 5G BS. Simulation results show that the proposed two-stage optimal dispatch method can effectively encourage multiple 5G BSs to participate in DR and achieve the win-win effect of assisting the ADN peak-shaving and low-carbon economic operation of 5G BSs.

Keywords: multiple PV-integrated 5G BSs, active distribution network, demand response, Lyapunov optimization, energy sharing

1 INTRODUCTION

The explosive growth of mobile data and the popularization of smart devices have accelerated the deployment of fifth-generation (5G) communication systems (Singh et al., 2020). However, while ensuring wide network coverage and high communication service quality, the high-power consumption characteristic of 5G base stations (BSs) not only imposes high electricity bills for communication operators but also exacerbates the non-negligible carbon emission problem (Piovesan et al., 2019). On the one hand, 5G BSs are equipped with 64/32-channel massive

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multiple-input multiple-output (MIMO) to realize high bandwidth and high communication traffic (Fu et al., 2019; Chih -Lin et al., 2020; Zhou et al., 2021), which causes the power consumption of a 5G BS to be 2–3 times higher than that of an ordinary 4G BS. On the other hand, the deployment density of 5G BSs is much higher than that of 4G BSs as 5G uses higher-frequency bands for communication (Han et al., 2021). Moreover, the standard 5G BS demands a power of more than 11.5kW (Israr et al., 2021), and the number of 5G BSs in China is expected to reach 13.1 million by 2025, with a total power consumption of 200 billion kWh. Therefore, the high-power consumption characteristic of 5G BSs has become the primary constraint to the ultra-dense network deployment in the 5G era, and it also impacts the stable operation of an active distribution network (ADN) (Pedram and Wang, 2019).

At present, powering BSs through distributed energy resources (DERs), such as photovoltaic (PV) generation and energy storage (ES), has become a common solution to reduce on-grid power consumption and build low-carbon wireless networks (Zhou et al., 2017; Hu et al., 2020). Although the introduction of DER devices increases the investment cost of communication operators, the cheapness and cleanliness of PV energy will greatly reduce the energy cost of BSs, and the spatio-temporal shifting characteristic of ES can further improve the absorption capacity of renewable energy. Moreover, 5G BSs are already equipped with battery backups during construction to ensure stable operation in case of power supply interruptions, which can be considered the inherent ES (Yong et al., 2021a; Tang et al., 2021). However, under the normal power supply status, large-scale distributed ESs in 5G BSs are always idle (Liao et al., 2020a; Ci et al., 2020). Therefore, on the basis of ensuring the uninterrupted power supply of the 5G BSs, if the energy source of the ES can be dispatched, the energy consumption flexibility of 5G BSs will be stimulated, and the high asset utilization of ES resources can be realized.

With the continuous transition from the traditional grid to smart grid (SG) and the widespread deployment of user-side smart meters (Khan and Jayaweera, 2019), PV-integrated 5G BSs participating in ADN demand response (DR) has become an irreplaceable way to reduce energy cost (Li et al., 2021). On the premise of ensuring the reliability of the 5G BS power supply, if the ES and PV resources in 5G BSs can be coordinated, more flexible dispatchable resources can be tapped to provide a powerful guarantee for the real-time balance of the source and load in the ADN. Moreover, communication operators can reduce electricity bills or obtain DR subsidy based on the time-of-use (TOU) tariff mechanism by fully exploiting the power utilization flexibility of PV-integrated 5G BSs and fully dispatching idle ES resources, thus reducing the energy cost of a 5G communication network (Yong et al., 2021b). Despite the advantages mentioned previously, the multiple PV-integrated 5G BSs participating in DR still confront several critical challenges, which are summarized as follows.

1) Because of the different spatio-temporal distributions of communication services (Huq et al., 2020) and light intensity, different degrees of power imbalance will occur in different PV-integrated 5G BSs, which increases the difficulty of 5G BS energy

dispatch in the ADN. In addition, as 5G BSs belong to communication operators, the direct energy dispatch of 5G BSs by the ADN is difficult to guarantee the communication service quality of telecom users and involves user security and privacy leakage issues.

2) The participation of multiple 5G BSs in the DR increases the operation and maintenance cost of ESs and impairs the participation enthusiasm of communication operators. Meanwhile, the power supply demands of 5G BSs with different spatio-temporal distributions vary greatly, and the randomness of PV energy and BS load further exacerbates the uncertainty of ES dispatchable capacity. How to effectively encourage and quantify the participation of 5G BSs in DR is an important problem to be solved.

3) During the participation of multiple 5G BSs in the DR, the imbalance between PV generation and load demand will lead to an over-discharge of ESs or the curtailment of PV energy. In addition, the accurate prediction of PV generation and load demand increases the difficulty of ensuring the real-time energy balance and puts higher requirements on real-time energy scheduling decisions (Wang et al., 2020). How to manage the energy of 5G BSs participating in DR is another challenge.

Multiple 5G BSs can be incentivized to participate in the DR through energy aggregators. It was pointed out that energy aggregators can aggregate BSs into a limited number of groups and act as third parties to dispatch BSs to participate in the DR (Xu et al., 2015). By introducing energy aggregators, encouraging and quantifying the participation of 5G BSs in the DR can be modeled as a social welfare maximization problem based on the contract theory, that is, maximizing the total utility of energy aggregators and BSs. A contract theory-based direct trading framework was proposed by Zhang et al. (2016) to address the lack of enthusiasm for direct trading between small-scale electricity suppliers and electricity consumers. Given the low enthusiasm of electric vehicles to participate in the DR, a contract theory-based charging rate allocation criterion and access control scheme were proposed by Zhang et al. (2018). However, the application of the contract theory can only ensure that multiple 5G BSs participate in the DR regularly and quantitatively in a large timescale but cannot guarantee the low-carbon and stable operation of 5G BSs in a small timescale.

The energy optimization of 5G BSs aims to alleviate the imbalance between PV generation and load demand by fully utilizing renewable energy, thereby reducing the energy cost and maintaining the stability of ES. An energy sharing model with energy aggregators as physical carriers was proposed by Guo et al. (2014), which allowed BSs to coordinate with each other for addressing energy imbalance by simultaneously transmitting and receiving energy to and from energy aggregators. Considering the unpredictability characteristics of the shared energy of 5G BSs, the Lyapunov optimization algorithm can be used to optimize the energy sharing mode, Liu et al. (2017) proposed an online energy management method based on Lyapunov optimization, aiming to fully absorb renewable energy and control the energy flow of the nano-grid group in real-time. Zhong et al. (2019) proposed an



online Lyapunov optimization-based control algorithm to optimize the shared energy in real-time in the multiuser-shared centralized ES system.

As the construction of 5G BSs is still in its infancy, research on the optimal energy dispatch of multiple PV-integrated 5G BSs participating in the DR is still relatively lacking. Motivated by the aforementioned challenges, we propose a two-stage optimal dispatch method for multiple PV-integrated 5G BSs participating in the ADN DR to fully dispatch PV and ES resources by combining Lyapunov optimization and contract theory. The main contributions of this study are summarized as follows.

- 1) The system architecture of multiple PV-integrated 5G BSs participating in the ADN DR is proposed, where the energy aggregator is introduced as the third party to motivate and assist the PV-integrated 5G BSs to participate in the DR.
- 2) A contract theory-based DR incentive mechanism is proposed to guide the discharge behavior of 5G BSs with different discharge capacities in the large-timescale DR planning stage. The peak-shaving demand of the ADN and the profit of BSs can be satisfied by optimizing the contract items.
- 3) A Lyapunov-based online energy optimization algorithm is proposed to make real-time decisions on the shared energy of multiple 5G BSs in the small timescale. The complex energy sharing problem is transformed into a linear programing problem, and only real-time information is required to make online decisions.

The remainder of this study is organized as follows. **Section 2** introduces the proposed system architecture and methodology of multiple PV-integrated 5G BSs participating in the ADN DR. **Section 3** elaborates the proposed contract theory-based large-timescale DR planning mechanism. In **Section 4**, the small-timescale online energy optimization algorithm for multiple

PV-integrated 5G BSs based on Lyapunov optimization is proposed. **Section 5** analyzes the numerical results, followed by the conclusion in **Section 6**.

2 SYSTEM ARCHITECTURE AND METHODOLOGY

2.1 System Architecture

The system architecture of multiple PV-integrated 5G BSs participating in the ADN DR is shown in Figure 1, which consists of a 5G communication network, an ADN, and an energy aggregator. The 5G communication network comprises I PV-integrated 5G BSs, which is denoted as $\mathcal{I} = \{BS_1, ..., BS_i, ..., BS_l\}$. Each 5G BS is equipped with an ES unit and a distributed PV generation device, denoted as $ES_i, \forall BS_i \in \mathcal{I}$ and $PV_i, \forall BS_i \in \mathcal{I}$, respectively. Because of the intermittent and volatile nature of PV energy, its high penetration in the ADN will lead to risk of overvoltage and continuous fluctuations in electricity prices. Therefore, we assume that the PV output from multiple PVintegrated 5G BSs can only interact with the energy aggregator to indirectly participate in the ADN DR. The ADN can provide a conventional power supply for the BSs and encourage the aggregator to provide auxiliary services such as frequency modulation and peakshaving during peak and valley load periods, but is unable to directly dispatch and control the energy resources of 5G BSs. Fortunately, the aggregator plays an intermediary role between the 5G communication network and the ADN and acts as a DR agent to encourage 5G BSs to participate in the DR and assist in energy sharing among 5G BSs. Under the coordination of the aggregator, the ADN and the 5G communication network are connected through power lines and communication links to support the interaction of energy flow and information flow (Lu et al., 2020).

- 1) Information flow interaction: in the DR planning stage, the incentive price signal of the aggregator and DR results of 5G BSs can be shared through two-way communication links. In the online energy optimization stage, the energy sharing requirements and results of 5G BSs also can be shared through two-way communication links.
- Energy flow interaction: with the assistance of the aggregator, 5G BSs can be used as response loads or sources. Specifically, 5G BSs can serve as typical prosumers of the ADN, which can both generate and consume electricity.



In this study, we explore how the aggregator encourages and assists multiple 5G BSs to participate in the ADN DR during peak load periods. Since the DR during peak load periods and the real-time energy management of 5G BSs are optimization problems on different timescales, the participation of 5G BSs in the DR includes two stages, that is, DR planning and online energy optimization. As shown in Figure 2, we utilize a multitimescale model (Yu et al., 2021) to divide the peak load period into T time slots, that is, small timescale, with length τ , the set of which is defined as $T = \{1, ..., t, ..., T\}$. Consecutive T_0 time slots are combined into a time interval, that is, large timescale, the set of which is defined as $M = \{1, ..., m, ..., M\}$. Therefore, the set of time slots contained in the m – th time interval can be denoted as $T(m) = \{(m-1)T_0 + 1, (m-1)T_0 + 2, ..., mT_0\}$. At the beginning of each interval, the aggregator designs a contract as an incentive mechanism to encourage the BSs to participate in the DR according to the discharge capacity of the 5G BSs. In each time slot, the aggregator performs the online energy optimization of 5G BSs based on the optimal contract determined in the large timescale.

2.2 Large-Timescale Demand Response Planning

In the DR planning stage, a contract theory-based largetimescale DR planning method is proposed, which takes the peak-shaving demand of the ADN as the constraint and maximizes social welfare to obtain the optimal discharge power of 5G BSs participating in the DR. Specifically, as the DR agent, the aggregator needs to coordinate multiple 5G BSs with discharge capacity to provide peak-shaving services during the peak load period. Because of the differences in power consumption, PV energy, and the state of charge (SoC) of ES, 5G BSs show different discharge capacities. Therefore, the aggregator will design different contract items for BSs with different discharge capacities to encourage them to respond to the peak-shaving demand of the ADN. Each contract item stipulates the relationship between the discharge power and reward. Each BS will select a contract item matching its type to provide stable power to obtain the best utility. The details of the contract design and optimization will be introduced in Section 3.

2.3 Small-Timescale Online Energy Optimization

In the online energy optimization stage, we focus on the impact of the real-time imbalance between the PV energy and load demand of BSs participating in the DR on the stable operation of ES. In the small timescale, the PV energy is first used to supply the load demand of the BS. When the PV energy is balanced with the load demand, the ES only needs to stably participate in the DR according to the contract item. When the PV energy is more than the load demand, surplus PV energy will be generated, which will affect the stable output of the ES and may cause PV energy abandonment. When the PV energy is less than the load demand, a load gap will be generated. To satisfy the communication reliability requirements, the ES needs to respond to the demands of the ADN while balancing the BS load gap, leading to over-discharge of ES. In addition, since the PV energy and load demand of BSs are usually independent, some BSs may lack sufficient PV energy to satisfy the load demand, while other BSs may abandon surplus PV energy because of the inability of complete absorption. Therefore, the aggregator is introduced as a shared energy carrier to support energy sharing among multiple 5G BSs, and a small-timescale online energy optimization algorithm is proposed to optimize the imbalance energy between the PV energy and load demand of multiple 5G BSs. The specific procedures will be introduced in **Section 4**.

3 CONTRACT THEORY-BASED DEMAND RESPONSE PLANNING

3.1 Base Station Type Modeling

In this study, we assume that the BS with more remaining ES energy and lower load demand has a larger discharge capacity and is more willing to participate in the DR to obtain higher rewards. For the aggregator, BSs can be classified into different types based on the discharge capacity. Assuming that in the m – th time interval, I 5G BSs can be classified into I BS types by arranging their discharge capacity in ascending order. The set of BS types can be denoted as $\Theta(m) = \{\theta_1(m), ..., \theta_i(m), ..., \theta_I(m)\}$, where $\theta_i(m)$ is discharge capacity BS_i the of and $\theta_1(m) \leq \cdots \leq \theta_i(m) \leq \cdots \leq \theta_I(m).$

The discharge capacity of a BS is determined by historical statistics of the source-load imbalance energy and ES energy in the first time slot of the m – th time interval (Wang et al., 2018). After the DR, the remaining ES energy should be greater than the minimum requirements of the ES discharge depth, which can be expressed as

$$E_i(t_0^m) + c_i - p_i(m) \cdot T_0 \tau \ge E_{\min}, t_0^m = (m-1)T_0 + 1, \quad (1)$$

where t_0^m is the first time slot of the m – th time interval. $E_i(t_0^m)$ is the initial ES energy in time slot t_0 . $T_0\tau$ is the duration of the time interval. $p_i(m)$ is the discharge power of BS_i in the m – th time interval. $p_i(m) \cdot T_0\tau$ is the total discharge electricity. c_i is the source-load imbalance margin, which is determined by the difference between the historical average PV energy and the load demand. When the historical average PV energy is greater than the load demand, c_i is a positive value. Otherwise, c_i is a negative value. E_{\min} is the lower bound of the remaining ES energy. According to (1), the discharge power of BS_i satisfies

$$p_i(m) \le \frac{E_i(t_0^m) - E_{\min} + c_i}{T_0 \tau}, t_0^m = (m-1)T_0 + 1.$$
 (2)

Therefore, the BS type of BS_i can be quantified as

$$\theta_i(m) = \frac{E_i(t_0^m) - E_{\min} + c_i}{T_0 \tau}, t_0^m = (m-1)T_0 + 1.$$
(3)

Considering the information security and privacy, the BS type of a specific BS is unknown to the aggregator, which means that the information is asymmetric (Zhou et al., 2019a). We assume that the aggregator can learn that there are a total of *I* types of BSs, and a BS belongs to type $\theta_i(m), \forall \theta_i(m) \in \Theta$ with probability P_i , that is, $\sum_{i=1}^{I} P_i = 1, \forall BS_i \in \mathcal{I}$.

3.2 Contract Formulation

The aggregator can design a contract including *I* different contract items for *I* types of BSs based on the contract theory to incentivize BSs to participate in the DR. For a type $\theta_i(m)$ BS, the aggregator pays a corresponding reward $r_i(m)$ for the discharge power $p_i(m)$ involved in the DR and specifies the performance-return correlation through the contract item $(p_i(m), r_i(m))$. The contract is denoted as $\pi = \{(p_i(m), r_i(m)), \forall BS_i \in \mathcal{I}\}.$

The principle of the aggregator contract design and BS contract signing is to maximize their own utility. Considering there are I types of 5G BSs, the utility of the aggregator is the subsidies for dispatching all response electricity minus all payment rewards, which is given by

$$U_{A}(\{p_{i}(m)\},\{r_{i}(m)\}) = I \sum_{i=1}^{N} P_{i}(T_{0}\tau\lambda(m)p_{i}(m) - r_{i}(m)), \quad (4)$$

where $\lambda(m)$ is the feed-in tariff of electricity in the m – th time interval. $T_0 \tau \lambda(m) p_i(m)$ is the DR subsidy received by the aggregator. The inequality $T_0 \tau \lambda(m) p_i(m) - r_i(m) \ge 0$ is always true for $\forall BS_i \in \mathcal{I}$, which means that the BSs participating in the DR are beneficial to the aggregator. Otherwise, the aggregator will lose momentum to aggregate energy of 5G BSs.

The utility of the type $\theta_i(m)$ BS which accepts the contract item $(p_i(m), r_i(m))$ is defined as the received reward minus the cost of discharging electricity, which is given by

$$U_{i}^{BS}(p_{i}(m), r_{i}(m)) = \theta_{i}(m)f(r_{i}(m)) - T_{0}\tau\xi p_{i}(m), \quad (5)$$

where ξ is the cost coefficient of discharging loss (Dragicevic et al., 2014). $\theta_i(m) f(r_i(m))$ is the value of $r_i(m)$ for type $\theta_i(m)$ BS, where the function $f(r_i(m))$ is a monotonically increasing concave function of $r_i(m)$ and represents the reward evaluation. Without the loss of generality, $f(r_i(m))$ is defined as a quadratic function, that is, $f(r_i(m)) = -ar_i^2(m)/2 + br_i(m)$, where *a* and *b* are positive constants which enable $f(r_i(m))$ to satisfy the constraints of f(0) = 0, $f'(r_i(m)) > 0$, and $f''(r_i(m)) < 0$.

According to the principle of contract design and signing, the problem of incentivizing 5G BSs to participate in the DR is transformed into the social welfare maximization problem. The expected social welfare is defined as the total utility of the aggregator and I BSs, which is given by

$$SW(\{p_i(m)\}, \{r_i(m)\}) = U_A(\{p_i(m)\}, \{r_i(m)\}) + I\sum_{i=1}^{I} P_i U_i^{BS}(p_i(m), r_i(m)).$$
(6)

3.3 Optimization Problem Modeling

The aggregator faces three constraints from the BS side when optimizing the contract design, namely, the individual rationality (IR) constraint, the incentive compatibility (IC) constraint, and the monotonicity constraint (Chen and Zhu, 2017). The IR constraint means that the BSs should receive non-negative utility after signing the contract and participating in the DR. The IC constraint means that a BS can only obtain maximum utility when signing the contract item designed for its own type. The monotonicity constraint means that the higher the BS type, the higher the reward will be. Therefore, the social welfare maximization problem is formulated as

P1:
$$\max_{\{(p_{i}(m),r_{i}(m))\}} SW(\{p_{i}(m)\},\{r_{i}(m)\})$$

s.t. C1: $\theta_{i}(m)f(r_{i}(m)) - T_{0}\tau\xi p_{i}(m) \ge 0$,
C2: $\theta_{i}(m)f(r_{i}(m)) - T_{0}\tau\xi p_{i}(m) \ge \theta_{i}(m)f(r_{i'}(m)) - T_{0}\tau\xi p_{i'}(m)$,
C3: $0 \le r_{1}(m) < \ldots < r_{i}(m) < \ldots < r_{I}(m)$,
C4: $p_{i}(m) \le \theta_{i}(m)$,
C5: $\sum_{i=1}^{I} p_{i}(m) = R(m) \quad \forall i, i' \in \{1, \dots, I\}$,
(7)

where C1, C2, and C3 are the IR constraint, IC constraint, and monotonicity constraint, respectively. C4 means that the discharge power demand stipulated in the contract should not exceed the discharge capacity of the BS. C5 is the peak-shaving demand constraint for power balance, where R(m) is the peak-shaving demand in the m-th time interval.

3.4 Problem Transformation and Optimal Contract Solution

The formulated optimization problem includes I IR constraints and I(I-1) IC constraints, the complexity of which is absolutely high. To facilitate the solution of the problem, the Contract Feasibility Necessary and Sufficient Condition Theorem (Zhou et al., 2019b) is used to reduce the dimensionality of IR constraints and IC constraints.

3.4.1 Individual Rationality Constraint Dimensionality Reduction

The utility of the type $\theta_i(m)$ ($i \neq 1$) BS should satisfy

$$\begin{aligned} \theta_i(m) f(r_i(m)) &- T_0 \tau \xi p_i(m) \ge \\ \theta_i(m) f(r_1(m)) &- T_0 \tau \xi p_1(m) \ge \\ \theta_1(m) f(r_1(m)) &- T_0 \tau \xi p_1(m) \ge 0, \end{aligned}$$

$$\end{aligned}$$

$$\tag{8}$$

where the first inequality satisfies the IC constraint, the second inequality satisfies the BS type definition, and the third inequality satisfies the IR constraint. Combining the first and third inequalities, we can draw that if the type $\theta_1(m)$ BS satisfies the IR constraint, the BSs with higher types will automatically satisfy the IR constraint. Therefore, only one IR constraint is retained through IR constraint dimensionality reduction, that is, $\theta_1(m) f(r_1(m)) - T_0 \tau \xi p_1(m) \ge 0$.

3.4.2 Incentive Compatibility Constraint Dimensionality Reduction

We consider three adjacent BS types, that is, $\theta_{i-1}(m) < \theta_i(m) < \theta_{i+1}(m)$, the utilities of which should satisfy



$$\theta_{i+1}(m) f(r_{i+1}(m)) - T_0 \tau \xi p_{i+1}(m) \ge \theta_{i+1}(m) f(r_i(m)) - T_0 \tau \xi p_i(m),$$
(9)

$$\theta_{i}(m) f(r_{i}(m)) - T_{0}\tau\xi p_{i}(m) \geq \theta_{i}(m) f(r_{i-1}(m)) - T_{0}\tau\xi p_{i-1}(m).$$
(10)

Based on $r_{i-1}(m) < r_i(m) < r_{i+1}(m)$, we have

$$\theta_{i+1}(m) f(r_{i+1}(m)) - T_0 \tau \xi p_{i+1}(m) \ge \theta_{i+1}(m) f(r_{i-1}(m)) - T_0 \tau \xi p_{i-1}(m).$$
(11)

Combining (9) and (11), it can be further concluded as

$$\theta_{i+1}(m)f(r_{i+1}(m)) - T_0\tau\xi p_{i+1}(m) \ge \theta_{i+1}(m)f(r_{i-1}(m))$$

$$-T_0\tau\xi p_{i-1}(m) \ge \cdots \ge \theta_{i+1}(m)f(r_1(m)) - T_0\tau\xi p_1(m).$$
(12)

constraint between the adjacent types holds. Through IC constraint dimensionality reduction, the number of IC constraints will be reduced from I(I-1) to I-1.

3.4.3 Optimal Contract Solution

Based on the aforementioned analysis, P1 is rewritten as follows:

P2:
$$\max_{\{(p_i(m),r_i(m))\}} SW(\{p_i(m)\},\{r_i(m)\})$$

s.t. C6: $\theta_1(m) f(r_1(m)) - T_0 \tau \xi p_1(m) \ge 0$,
C7: $\theta_i(m) f(r_{i-1}(m)) - T_0 \tau \xi p_{i-1}(m) \le \theta_i(m) f(r_i(m)) - T_0 \tau \xi p_i(m)$,
C3, C4, C5 $\forall i \in \{1, \cdots, I\}$.
(13)

By checking the Hessian matrix, it can be known that the objective function of P2 is concave. Because the difference between two concave functions is involved in C7, P2 cannot be solved directly by convex optimization. Therefore, we adopt the concave-convex process (CCP) algorithm (Wei et al., 2016) to transform P2 into a convex optimization problem. Denote $F_i(r_i(m)) = \theta_i(m) f(r_i(m))$, which can be differentiable with regards to $r_i(m)$. Then, $F_i(r_i(m))$ can be approximated by the first-order Taylor series expansion as

$$F_i(r_i(m)) \approx F_i(r_{i,0}[s]) + \nabla F_i(r_{i,0}[s])(r_i(m) - r_{i,0}[s]), \quad (14)$$

where $r_{i,0}[s]$ represents the initial point in the s – th iteration. Therefore, *C*7 is converted to the difference between a concave function and an affine function, which is expressed as

$$\tilde{C}7: \theta_{i}(m)f(r_{i-1}(m)) - T_{0}\tau\xi p_{i-1}(m) \\\leq \theta_{i}(m) \left[F_{i}(r_{i,0}[s]) + \nabla F_{i}(r_{i,0}[s])(r_{i}(m) - r_{i,0}[s])\right] - T_{0}\tau\xi p_{i}(m).$$
(15)

By replacing C7 with $\tilde{C}7$, P2 is transformed into a convex optimization problem. In the *s* – th iteration, the local optimal solutions $\hat{r}_i[s]$ and $\hat{p}_i[s]$ can be obtained by solving the transformed convex optimization problem. Then, the initial point of the Taylor series expansion in the (s + 1) – th iteration is defined as $r_{i,0}[s+1] = \hat{r}_i[s]$, and the new local optimal solutions are derived until the iteration stop condition is met, that is, the change in the expected social welfare is less than or equal to a certain threshold, which is given by

$$SW(\{\hat{p}_{i}[s+1]\}, \{\hat{r}_{i}[s+1]\}) - SW(\{\hat{p}_{i}[s]\}, \{\hat{r}_{i}[s]\}) \le \varepsilon.$$
(16)

Then, the iteration is terminated and the contract design is completed. The discharge power and corresponding reward of 5G BSs participating in the DR can be obtained based on the contract.

4 LYAPUNOV-BASED ONLINE ENERGY OPTIMIZATION

4.1 Dynamic Energy Queue Model

On the basis of obtaining the optimal discharge power of 5G BSs participating in the DR, we analyze the energy flow of BSs in the small timescale and propose the energy sharing strategy among multiple 5G BSs to further reduce the energy cost of the 5G communication network. Specifically, we classify the BSs into energy output BSs and energy input BSs, where energy output BSs share energy to absorb the surplus PV energy and energy input BSs receive the shared energy to balance the load gap. We define \mathcal{BS}^+ and \mathcal{BS}^- as the sets of energy output BSs and energy input BSs, respectively. The load demand and the PV energy of BS_i in time slot t are denoted as $L_i(t), t \in T(m)$ and $S_i(t), t \in T(m)$, respectively. The surplus PV energy of BS_i in time slot t is denoted as $C_i(t), t \in T(m)$, which is given by

$$C_{i}(t) = \begin{cases} \min(S_{i}(t) - L_{i}(t), C_{\max}), BS_{i} \in \mathcal{BS}^{+}\\ 0, \qquad BS_{i} \in \mathcal{BS}^{-}_{i} \end{cases}, \quad (17)$$

where C_{max} is the maximum surplus PV energy limited by the transmission line capacity, and the energy sharing among BSs obeys the following principles.



TABLE 1 | Simulation parameters.

Project	Attribute	Value	
PV generation	Capacity/kWp	10	
	Power generation/kWh	64.8227	
ES unit	Capacity/kWh	20.24	
	Maximum discharging power/kW	10.24	
	Lowest level of SoC/%	30	
	Cost coefficient of discharge loss	0.14	
ES initial SoC/%	BS1	38	
	BS2	39	
	BS3	43	
	BS4	46	
	BS5	48	
	BS6	52	
	BS7	55	
	BS8	56	
	BS9	59	
	BS10	61	

- 1) For $BS_i \in \mathcal{BS}^+$, the shared energy only comes from $C_i(t)$, and ES_i does not participate in energy sharing. When there is no load gap in other BSs, $C_i(t)$ and ES_i will participate in the DR together.
- For BS_i ∈ BS⁻, the input energy can only be used to balance the load gap to reduce the discharge depth of ES_i and cannot participate in the DR.

Therefore, in the process of energy sharing, $BS_i \in \mathcal{BS}^+$ can share energy to other energy input BSs for absorbing surplus PV energy, and $BS_i \in \mathcal{BS}^-$ can balance the load gap by obtaining input energy to reduce the excessive utilization of ES_i energy. It is to be noted that when the input energy is insufficient to balance the load gap, the ES_i will increase the discharge energy to satisfy the load demand for $BS_i \in \mathcal{BS}^-$. The energy flows of \mathcal{BS}^+ and \mathcal{BS}^- are shown in **Figure 3**.

The set of shared energy of BSs in time slot t can be defined as $X(t) = [X_1(t), X_2(t), \dots, X_I(t)]$, where $X_i(t)$ is the shared

energy of BS_i . Denote $X_i(t)$ as $X_i(t)^+$ if $X_i(t) > 0$, which means the surplus PV energy $C_i(t) > 0$. Denote $X_i(t)$ as $X_i(t)^-$ if $X_i(t) < 0$, which means BS_i confronts the load gap. The power balance constraint and upper/lower bound constraint of $X_i(t)$ can be expressed as

$$\sum_{i} X_{i}(t) = 0 \text{ and } 0 \le |X_{i}(t)| \le X_{\max},$$
(18)

where X_{max} is the maximum transmission energy of the shared energy. In addition, based on the basic principle of energy sharing, the shared energy should satisfy

$$X_{i}(t) = \begin{cases} 0 \le X_{i}(t)^{+} \le C_{i}(t), & BS_{i} \in \mathcal{BS}^{+} \\ -(L_{i}(t) - S_{i}(t)) \le X_{i}(t)^{-} < 0, BS_{i} \in \mathcal{BS}^{-} \end{cases}.$$
(19)

Without the energy sharing mode, the set of discharge energy of ESs in time slot *t* is denoted as $B(t) = [B_1(t), B_2(t), \dots, B_I(t)]$. The discharge energy of ES_i in time slot *t* is given by

$$B_{i}(t) = \begin{cases} p_{i}(m) \cdot \tau, & BS_{i} \in \mathcal{BS}^{+} \\ \min\left(\left[p_{i}(m) \cdot \tau + L_{i}(t) - S_{i}(t)\right], B_{\max}\right), BS_{i} \in \mathcal{BS}^{-}, \end{cases}$$

$$(20)$$

where B_{max} is the maximum discharge energy, which depends on the maximum discharge power of ES.

The energy level of ES can be regarded as an ES queue, the length of which represents the remaining energy of the ES. We define the set of ES queues in time slot *t* as $E(t) = [E_1(t), E_2(t), \dots, E_I(t)]$, where $E_{\min} \leq E_i(t) \leq E_{\max}$. According to the energy flow of \mathcal{BS}^+ and \mathcal{BS}^- , the dynamic update equation of $E_i(t)$ can be further obtained as

$$E_{i}(t+1) = \begin{cases} E_{i}(t) - (B_{i}(t) + X_{i}(t)^{+} - C_{i}(t)), & BS_{i} \in \mathcal{BS}^{+} \\ E_{i}(t) - (p_{i}(m) \cdot \tau + L_{i}(t) - S_{i}(t) + X_{i}(t)^{-}), & BS_{i} \in \mathcal{BS}^{-} \end{cases}$$
(21)

E(t) is called to be mean rate stable (Zhou et al., 2019c) if it satisfies

$$\lim_{t \to +\infty} \frac{\mathbb{E}\{|E_i(t)|\}}{\tau} = 0, \forall i \in \{1, \cdots, I\}.$$
(22)

The reason for ensuring the stability of ES queues during 5G BSs participating in the DR is that the frequent and excessive discharge of ESs will seriously shorten the life span and increase the maintenance cost of ES.

4.2 Online Energy Optimization Problem Formulation

The objective of the online energy optimization problem is to realize the full absorption of PV energy by energy sharing among 5G BSs, while ensuring the stability of ES queues. The PV absorption rate is denoted as the percentage of PV energy that can be absorbed by 5G BSs, which is given by

$$K_{pcr}(t) = \frac{\sum_{i} (X_{i}(t)^{+} + N_{i}(t))}{\sum_{i} L_{i}(t)},$$
(23)

where $N_i(t)$ is the self-produced PV energy $S_i(t)$ absorbed by the load demand of BS_i in time slot *t*, and $N_i(t) = \min(S_i(t), L_i(t))$.

TABLE 2 | Optimal contract parameters.

5G BS	Discharge capacity/kW	Response power/kW	BS utility	Aggregator utility	Social welfare
1	1.5	1.5	0.67	0.45	1.12
2	2.0	2.0	1.01	0.64	1.65
3	2.5	2.5	1.37	0.84	2.21
4	3.0	3.0	1.75	1.04	2.79
5	3.5	3.5	2.15	1.24	3.39
6	4.0	4.0	2.55	1.45	4.00
7	4.5	4.5	2.97	1.65	4.68
8	5.0	5.0	3.40	1.85	5.25
9	5.5	5.5	3.84	2.06	5.90
10	6.0	6.0	4.29	2.26	6.55
Summation	37.5	37.5	24	13.48	37.54



Furthermore, the time-average PV absorption rate can be described as

$$\bar{K}_{pcr}(t) = \lim_{T_0 \to \infty} \frac{1}{T_0} \sum_{t=(m-1)T_0+1}^{mT_0} \frac{\sum_i (X_i(t)^+ + N_i(t))}{\sum_i L_i(t)}.$$
 (24)

The focus of online energy optimization is formulated to improve the PV absorption rate during the whole peak load period to minimize the energy cost of multiple 5G BSs participating in the DR by real-time decision-making on the shared energy among BSs, which can be expressed as

P3:
$$\max_{X_i(t)^+} \bar{K}_{pcr}(t)$$

s.t. (18), (19), (21), (22), $\forall i, t.$ (25)

As P3 is a time-coupling optimization problem due to the long-term constraints, traditional methods, such as stochastic optimization and dynamic programming, will suffer from the problem of dimensionality. Lyapunov optimization is an effective method to solve long-term optimization problems, which can transform the coupled long-term optimization problem into an independent single-slot deterministic sub-problem. Compared with traditional methods, Lyapunov optimization requires less prior information and poses lower computational complexity (Liao et al., 2020b). Therefore, we propose a Lyapunov-based online energy optimization algorithm for multiple 5G BSs participating in the DR.

4.3 Lyapunov-Based Problem Transformation and Online Energy Optimization Algorithm

The Lyapunov function is defined as $L(\boldsymbol{E}(t)) = \sum_i E_i(t)^2/2$, which represents a scalar measure of the remaining energy in ES queues. To reflect the stability of the ES queues, we define the one-slot conditional Lyapunov drift as the conditional expected change of the Lyapunov function between two adjacent time slots, which is given by

$$\Delta(\boldsymbol{E}(t)) = \mathbb{E}\{L(\boldsymbol{E}(t+1)) - L(\boldsymbol{E}(t))|\boldsymbol{E}(t)\}.$$
(26)

As (25) is a maximization problem, we define the drift-minusreward (Li et al., 2020) to maximize the PV absorption rate while ensuring the queue stability, which is given by

	45		I	I	T	Т	T	1	T	1	T		BS1	
	30		1		1			1	1		1	-		-
	0		5	10	15	20	25	30	35	40	45	50	55	60
	45		1	1	1	1	1	1	1	1	1		BS2	
	30			· · · · ·	· · · · ·			1	1					
	0		5	10	15	20	25	30	35	40	45	50	55	60
	50		I	1	T	1	1	1		1	Т		BS3	
	30				-1									
	0		5	10	15	20	25	30	35	40	45	50	55	60
	50						I	1	1	1	T		BS4	
	30													-
	0		5	10	15	20	25	30	35	40	45	50	55	60
	55				1	1	1	1	1	1	1		BS5	
%	30		·											
oC/o	0		5	10	15	20	25	30	35	40	45	50	55	60
S	55		-				Ι	1	I		1		BS6	
	30		F											
	0		5	10	15	20	25	30	35	40	45	50	55	60
	60				1	1	1	1	1	I	1		BS7	
	30		ł		-i									
	0		5	10	15	20	25	30	35	40	45	50	55	60
	60				1	1	I	Т	1	1	T		BS8	
	30		ł		- i		<u>j</u>							
	0		5	10	15	20	25	30	35	40	45	50	55	60
	65					1	I	1	1	I	1		BS9	
	30		+											
	0		5	10	15	20	25	30	35	40	45	50	55	60
	65									I	1		BS10	
	30		·											
	0		5	10	15	20	25 M	30 linute	35	40	45	50	55	60
	Without Energy Sharing Energy Sharing SoC Minimum Limit								- +					
	į													<u> </u>
FI	FIGURE 6 Comparison of the ES SoC.													

P4:
$$\min_{X_i(t)^+, X_i(t)^-} \Delta(\boldsymbol{E}(t)) - V\mathbb{E}\left\{K_{pcr}(t)\right\}$$

s.t. (18), (19), (21), $\forall i, t,$ (27)

where V is a non-negative weight parameter used to make a tradeoff between ES queue stability and PV absorption rate maximization. In the case of all possible E(t) and $V \ge 0$, the drift-minus-reward is upper-bounded by

$$\Delta(\boldsymbol{E}(t)) - \mathbb{E}\left\{V \cdot \frac{\sum_{i} (X_{i}(t)^{+} + N_{i}(t))}{\sum_{i} L_{i}(t)}\right\} \\
\leq \sum_{i} \mathbb{E}\{E_{i}(t) [C_{i}(t) - X_{i}(t)^{+} - B_{i}(t)] | X_{i}(t) \ge 0\} \\
+ \sum_{i} \mathbb{E}\{E_{i}(t) [-p_{i}(m) \cdot \tau - L_{i}(t) + S_{i}(t) - X_{i}(t)^{-}] | X_{i}(t) < 0\} \\
- \mathbb{E}\left\{V \cdot \frac{\sum_{i} (X_{i}(t)^{+} + N_{i}(t))}{\sum_{i} L_{i}(t)}\right\} + Z,$$
(28)

where Z is a positive constant and is given by

$$Z = \sum_{i} \left(B_{\max}^{2} + X_{\max}^{2} \right).$$
 (29)

Proof: see Supplementary material.

Therefore, P4 can be further transformed into a linear programing problem by taking the iterative expectation and simplifying the upper bound of $\Delta(E(t)) - V\mathbb{E}\{K_{pcr}(t)\}$ derived in (28), which is given by

P5:
$$\min_{X_{i}(t)^{*},X_{i}(t)^{-}} - \frac{V}{\sum_{i} L_{i}(t)} \sum_{i} X_{i}(t)^{*} - \sum_{i} E_{i}(t) X_{i}(t)^{*} - \sum_{i} E_{i}(t) X_{i}(t)^{-}$$

s.t. (18), (19), (21), $\forall i, t,$ (30)

Since $E_i(t)$ and $L_i(t)$ are available in time slot t, P5 is a linear programing problem with shared energy as the optimization variable, which can be solved directly using the CPLEX solver. It is observed that the Lyapunov-based online energy optimization algorithm only needs the real-time information of the remaining energy of ESs and the load demand of BSs to make the energy sharing decision, which effectively improves the efficiency of the algorithm execution. Moreover, the computational complexity of the proposed online energy optimization algorithm only increases linearly with the expansion of the number of BSs.



TABLE 3 | Comparison of the PV absorption rate of 5G BSs.

and the second							
Indicator	Without DR participation	DR participation	Improvement/%				
Abandoned PV/kWh	11.41	0.086	17.47				
PV absorption rate/%	82.4	99.87					

5 CASE STUDY

5.1 Basic Data

In this section, we validate the performance of the proposed twostage optimal dispatch method by simulations. We consider a scenario in which the aggregator signs the DR contract with 10 5G BSs during the peak load period. The BSs provide peak-shaving service for the ADN from 10:00 a.m. to 11:00 a.m. The feed-in tariff is CNY 0.42/kWh, and the peak-shaving demand is 37.5 kW. The case analysis uses 1 min as the small timescale and 1 h as the large timescale. The minimum load demand is 2.193 kW when the BS is in a no-load operation state, and the maximum load demand is 10 kW when the BS is in a full-load operation state (China Mobile Research Institute, 2020). Considering that 10 5G BSs are geographically close to each other, we assume that the PV output of each BS adopts the same data. The typical PV output curve of 5G BSs from 10:00 a.m. to 11:00 a.m. is shown in Figure 4, and the simulation parameter settings of the DERs in 5G BSs are shown in Table 1. Moreover, since the daily dynamic load data of 5G BSs are still undisclosed, we assume that the load demand of 10 5G BSs obeys the random distribution within the load demand interval (2.193 kW, 10 kW).

5.2 Comparative Analysis of Optimization Results

5.2.1 Contract Theory–Based Large-Timescale Demand Response Planning

According to the initial ES capacity of the 5G BSs, the minutelevel power consumption, and the PV output curve, the optimal contract can be obtained, and the corresponding parameters are shown in **Table 2**. The results show that the response power coincides with the discharge capacity, which implies that the 5G BSs with different discharge capacities choose the contract items that conform to their own types. The optimal discharge power is provided to obtain the optimal reward. The utilities of BSs and the aggregator increase with the BS type, which indicates that the aggregator prefers to sign the contract with BSs with higher discharge capacity. Moreover, we can find that the sum of the response power is equal to 37.5 kW, which accurately satisfies the peak-shaving demand of the ADN.

The optimal discharge power and discharge reward versus the BS type are shown in **Figure 5A**. It can be seen that the optimal discharge power and discharge reward increase monotonously with the BS type, indicating that the contract satisfies the monotonicity constraint, and the discharge capabilities of BSs are well-exploited. **Figure 5B** shows the utility of 5G BSs versus the type of the contract item. It can be seen that maximum utility can be achieved only when the 5G BS signs the contract item

conforming to its type, indicating that the contract satisfies the IC constraint. Moreover, all 5G BSs that sign the contract can obtain non-negative utility, which is consistent with the IR constraint. In addition, the utility of the 5G BS is proportional to the BS type, which will increase the motivation of the higher type of 5G BS to sign the contract and participate in the DR.

5.2.2 Lyapunov-Based Small-Timescale Online Energy Optimization

Figure 6 shows the comparison of ES SoC with and without energy sharing. Simulation results show that 5G BSs participating in the DR without energy sharing will lead to an over-discharge or insufficient dispatching of ESs. At the end of DR, the SoCs of BS1 and BS2 are lower than the minimum capacity limit, while those of BS6 and BS7 still have enough dispatchable space. By contrast, the proposed Lyapunov-based online energy optimization algorithm with energy sharing can achieve higher discharge stability and a more balanced power level of ESs.

Figure 7 shows the shared energy and source-load imbalance energy of 10 5G BSs during the DR period. The red curve and the blue cylinder represent the shared energy optimized by the proposed online energy optimization algorithm and the imbalance energy, respectively. When the red curve is tangential to the blue cylinder, it means that the shared energy completely compensates for the imbalance energy. At this point, the ES unit is in a stable output state and only needs to output constant power based on the contract item. From 10:00 a.m. to 11:00 a.m., most BSs have surplus PV energy due to sufficient sunlight, and only a few BSs with high power consumption require shared energy. Based on the proposed online energy optimization algorithm, it can be seen that most BSs can interact with each other to compensate for the imbalance in energy through energy sharing, which further promotes PV energy consumption, reduces the curtailment of PV output, and ensures the output stability of ES units.

Table 3 shows the comparison of the PV absorption rate before and after the two-stage optimal dispatch. Compared with not participating in the DR, 5G BSs participating in the DR can improve the PV absorption rate by 17.47% and make it up to 99.87%. The reason is that the contract-based DR planning ensures rough absorption of the overall PV output during the entire peak load period, while the Lyapunov-based online energy optimization algorithm ensures the approximately complete and real-time absorption of the PV output, thus realizing the maximum absorption of PV energy during the entire peak load period.

6 CONCLUSION

In this study, we propose a two-stage optimal dispatch method based on the contract theory and Lyapunov optimization for the typical scenario of multiple PV-integrated 5G BSs participating in the ADN DR. The following conclusions are obtained:

- 1) The system architecture for multiple PV-integrated 5G BS participating in the ADN DR was proposed, which is composed of aggregators, a 5G communication network, and an ADN.
- 2) A contract theory-based large-timescale DR planning method was proposed for modeling the interaction between 5G BSs and the aggregator. Simulation results show that the proposed method can improve the motivation of 5G BSs to participate in the DR and satisfy the peak-shaving demand of the ADN.
- 3) A Lyapunov-based small-timescale online energy optimization algorithm was proposed for making realtime decisions on energy sharing of 5G BSs. The numerical results show that the proposed method can improve the PV absorption rate up to 99.87%, which achieves approximately complete absorption of PV energy and ensures the stability of the ES queue.

Further research will be conducted in three directions, including:

1) Implementation of energy dispatch for large-scale 5G BSs at multiple timescales including day-ahead, intra-day, and realtime and the systematical evaluation of the respond capacity of a PV-integrated 5G BS for ancillary services.

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- 2) Energy sharing scheme and a demand response incentive mechanism that take into account power transmission losses.
- 3) Multi-entity benefit allocation method for participating in the DR under deregulated power markets.

DATA AVAILABILITY STATEMENT

The original contributions presented in the study are included in the article/Supplementary Material; further inquiries can be directed to the corresponding author.

AUTHOR CONTRIBUTIONS

XZ: conceptualization, methodology, software, investigation, formal analysis, and writing-original draft; ZAW: data curation and writing-original draft; ZZ: visualization and writing-original draft; HL: conceptualization, supervision, and writing-review and editing; XM: conceptualization, supervision, and writing-review and editing; XY: visualization and writing-review and editing GL: resources, supervision, and writing-review and editing; ZOW: visualization and editing; writing-review and ZL: visualization and writing-review and editing; YL: resources and supervision.

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APPENDIX

Proof of (27):

$$\Delta(E(t)) = \mathbb{E}\{L(E(t+1)) - L(E(t))\}$$

$$= \frac{1}{2}\mathbb{E}\{\sum_{i} [E_{i}(t+1)^{2} - E_{i}(t)^{2}]\}$$

$$= \frac{1}{2}\mathbb{E}\{\sum_{i} [(E_{i}(t) - B_{i}(t) + C_{i}(t) - X_{i}(t)^{+})^{2} - E_{i}(t)^{2}]|X_{i}(t) \ge 0\}$$

$$+ \frac{1}{2}\mathbb{E}\{\sum_{i} [(E_{i}(t) - p_{i}(m) \cdot \tau - (L_{i}(t) - S_{i}(t)) - X_{i}(t)^{-})^{2} - E_{i}(t)^{2}]|X_{i}(t) < 0\}$$

$$\leq \frac{1}{2}\mathbb{E}\{\sum_{i} [B_{i}(t)^{2} + C_{i}(t)^{2} + (X_{i}(t)^{+})^{2} - 2E_{i}(t) (C_{i}(t) - X_{i}(t)^{+} - B_{i}(t))]|X_{i}(t) \ge 0\} + \frac{1}{2}\mathbb{E}\{\sum_{i} [p_{i}(m)^{2} \cdot \tau^{2} + (L_{i}(t) - S_{i}(t))^{2} + (X_{i}(t)^{-})^{2}$$

Thus, we can have (28)

$$\begin{split} \Delta E(t) &- \mathbb{E}\left\{ V \cdot \frac{\sum_{i} (X_{i}(t)^{+} + N_{i}(t))}{\sum_{i} L_{i}(t)} \right\} \\ &\leq \sum_{i} \mathbb{E}\{E_{i}(t) [C_{i}(t) - X_{i}(t)^{+} - B_{i}(t)] | X_{i}(t) \ge 0\} \\ &+ \sum_{i} \mathbb{E}\{E_{i}(t) [-p_{i}(m) \cdot \tau - L_{i}(t) + S_{i}(t) - X_{i}(t)^{-}] | X_{i}(t) < 0\} \\ &- \mathbb{E}\left\{ V \cdot \frac{\sum_{i} (X_{i}(t)^{+} + N_{i}(t))}{\sum_{i} L_{i}(t)} \right\} + Z \end{split}$$

The proof is completed.